

A note on the algebra of threshold graphs

Research Article

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Abstract: It is known that threshold graphs have edge rings with 2-linear resolutions. This was proved by Engström and Stamps, [4]. They used the fact that an edge ring of a graph G has a 2-linear resolution if and only if the complement graph is chordal. They also described a method to determine the Betti numbers. Our goal is to determine when edge rings of threshold graphs are Cohen-Macaulay. In order to do so, it is more convenient to use an alternative way to study edge rings of graphs, that is to interpret them as Stanley-Reisner rings. We also determine when the neighborhood complex of a threshold graph has a Cohen-Macaulay Stanley-Reisner ring.

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1. Introduction

A simplicial complex Δ consists of a set of subsets (faces) σ of a finite set $V = \{v_1, \dots, v_n\}$, such that $\{v_i\} \in \Delta$ for all i , and if $\sigma_1 \in \Delta$ and $\sigma_2 \subset \sigma_1$, then $\sigma_2 \in \Delta$. Maximal faces are called facets. If Δ is a simplicial complex with vertices $\{v_1, \dots, v_n\}$, its Stanley-Reisner ring over a field k is $k[\Delta] = k[x_1, \dots, x_n]/I_\Delta$, where its Stanley-Reisner ideal I_Δ is generated by the squarefree monomials $x_{i_1} \cdots x_{i_k}$ for all non-faces $\{v_{i_1}, \dots, v_{i_k}\}$ of Δ . If Δ is a simplicial complex on $\{1, 2, \dots, n\} = [1, n]$, then the Alexander dual Δ^\vee is the complex on $[1, n]$ for which $\sigma \subseteq [1, n]$ is a face of Δ^\vee if and only if $[1, n] \setminus \sigma$ is not a face of Δ . The f -vector of a simplicial complex Δ on $[1, n]$ is $f(\Delta) = (f_{-1}, f_0, \dots, f_{n-1})$, where f_i is the number of i -dimensional faces in Δ . Let $f_{-1} = 1$. We usually skip the zeros at the end.

A graph $G = (V, E)$ is a finite set of vertices $V = \{v_1, \dots, v_n\}$ and edges $E = \{(v_i, v_j)\}$. A vertex v is called isolated if v is not contained in any edge. A subset $W \subseteq V$ is called independent if the graph induced on W consists of isolated vertices. Two vertices v and w are called neighbors if (v, w) is an edge. For a vertex v , let $N(v)$ consist of all neighbors to v . The edge ring of $G = (V, E)$ is $k[G] = k[x_1, \dots, x_n]/I(G)$, where $V = \{v_1, \dots, v_n\}$ and the edge ideal $I(G)$ is generated by $\{x_i x_j; (v_i, v_j) \in E\}$. The complement graph to $G = (V, E)$ is $\bar{G} = (V, \bar{E})$, where $\bar{E} = \{(x_i, x_j); (i, j) \notin E\}$.

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A graded algebra $R = k[x_1, \dots, x_n]/I = S/I = \bigoplus_{i \geq 0} R_i$ has Hilbert series $R(t) = \sum_{i \geq 0} \dim_k R_i t^i$. It has a t -linear resolution if I is generated in degree t and all higher syzygies are linear, or equivalently, if $\text{Tor}_{i,j}^S(R, k) = 0$ if $j \neq i + t - 1$. A t -linear algebra R has Hilbert series $R(t) = (1 + \sum_{i=1}^n (-1)^i \beta_{i,i+t-1}(R) t^{i+t-1}) / (1-t)^n$, where $\beta_{i,j}(R) = \dim_k \text{Tor}_{i,j}^S(R, k)$. Thus, if R has a t -linear resolution, the Hilbert series contains the same information as the set of all Betti numbers. We will concentrate on 2-linear resolutions and denote the vector $(\beta_{1,2}(R), \beta_{2,3}(R), \dots, \beta_{n,n+1}(R))$ by $\beta(R)$. We skip zeros at the end.

We now define some graph complexes, i.e., simplicial complexes belonging to a graph.

Definition 1.1. For a graph $G = (V, E)$, the neighborhood complex $N(G)$ consists of all subsets of $N(v)$ for all $v \in V$.

The dominance complex $D(G)$ of a graph $G = (V, E)$ consists of those $W \subseteq V$ such that $V \setminus W$ is dominating. A subset U of V is dominating if every vertex in V either belongs to U or has a neighbor in U .

The independence complex $\text{Ind}(G)$ of a graph $G = (V, E)$ consists of the independent subsets of V . The Stanley-Reisner ideal $I_{\text{Ind}(G)}$ equals the edge ideal $I(G)$.

The definition of threshold graphs, which also explains the name, stems from [1].

Definition 1.2. (Chvátal and Hammer, [1]) A graph $G = (V, E)$ is called a threshold graph when there exist nonnegative reals (or integers) w_v and t such that $\sum_{v \in U} w_v \leq t$ if and only if U is an independent subset of V .

Theorem 1.3. The following conditions are equivalent for a graph G :

1. G is a threshold graph.
2. G is a split graph, i.e., V can be partitioned in a clique and an independent set, and the neighborhoods of vertices in the independent set are nested.
3. G can be constructed from a vertex by repeatedly adding an isolated vertex or a star vertex (i.e. with edges to all previous vertices).
4. The dominance complex $D(G)$ of G equals the neighborhood complex $N(G)$ of G .

Proof. The equivalence of 1., 2. and 3. is proved in [8]. The equivalence of 1. and 4. is in [7]. □

There are more equivalent conditions proved in [8] and [7]. Here we add a further equivalent condition.

Theorem 1.4. The Alexander dual to the neighborhood complex of G equals the neighborhood complex of the complement to G , that is $N(G)^\vee = N(\overline{G})$, if and only if G is threshold.

Proof. That for any G we have $D(G) = N(\overline{G})^\vee$ is proved in [9]. Thus according to 4., we get that 1. is equivalent to $N(\overline{G})^\vee = N(G)$, and so is $N(\overline{G}) = N(G)^\vee$. □

Theorem 1.5. [5] For a graph $G = ([1, n], E)$, $k[x_1, \dots, x_n]/I(G)$ has a 2-linear resolution if and only if \overline{G} is a chordal graph, i.e., every cycle in \overline{G} of length > 3 has a chord.

This theorem is made more precise in [2]. If $W \subset V$ let $G[W]$ be the graph induced on W .

Theorem 1.6. [2, Theorem 3.2] Let \overline{G} be a chordal graph. If $i \neq j - 1$ then $\beta_{i,j}(k[G]) = 0$ and otherwise $\beta_{i,j}(k[G]) = \sum_{W \in \binom{V}{j}} (-1 + \# \text{ connected components of } \overline{G}[W])$.

It is also known which simplicial complexes that have a Stanley-Reisner ring with 2-linear resolution. These are the so called fat forests. Here is a recursive definition of fat forests, [5], [6]. A simplex F_1 is a fat forest. If $F_i, i = 1, \dots, k$ are simplices and $G_{k-1} = \cup_{i=1}^{k-1} F_i$ is a fat forest, then $G_{k-1} \cup F_k$ is a fat forest if $H = G_{k-1} \cap F_k$ is a simplex, $\dim H \geq -1$. (If $\dim H = -1$, then G_{k-1} and F_k are disjoint.) As an example, let F_1, \dots, F_4 be the simplices on $\{1, 2, 3, 4\}, \{1, 2, 3, 5\}, \{1, 2, 6\}, \{5, 7\}$ respectively. The intersections are the simplices on $\{1, 2, 3\}, \{1, 2\}$, and $\{5\}$, respectively. Then we build the fat forest with vertices $\{1, 2, 3, 4, 5, 6, 7\}$ and maximal faces F_1, F_2, F_3, F_4 .

Theorem 1.7. [5, 6] $k[\Delta]$ has a 2-linear resolution if and only if Δ is a fat forest.

The projective dimension of a ring $R = S/I$ is $\max\{i; \text{Tor}_{i,j}^S \neq 0\}$. The regularity of R is $\max\{j - i; \text{Tor}_{i,j}^S \neq 0\}$. The depth of R is the maximal length of a regular sequence in R . Finally, R is Cohen-Macaulay if $\text{depth}(R) = \dim(R)$. It is easy to get the Hilbert series, the dimension, and the depth for the Stanley-Reisner ring of a fat forest.

Theorem 1.8. [6, Theorem 1] Let $F = F_1 \cup \dots \cup F_k$ be a fat forest with F_i a simplex of dimension d_i and $F_1 \cup \dots \cup F_{j-1} \cap F_j$ a simplex of dimension r_j . Then the Hilbert series of $k[F]$ is

$$\sum_{i=1}^k \frac{1}{(1-t)^{d_i+1}} - \sum_{i=2}^k \frac{1}{(1-t)^{r_i+1}}.$$

The dimension of $k[G]$ is $\max\{d_i + 1\}_{i=1}^k$ and the depth of $k[G]$ is $\min\{r_i + 1\}_{i=1}^k$. $k[G]$ is Cohen-Macaulay if and only if there is a d such that $d_i = r_i + 1 = d$ for all i .

Theorem 1.9. [3, Theorem 3][10, 11] $k[\Sigma]$ has a linear resolution if and only if $k[\Sigma^\vee]$ is Cohen-Macaulay. The regularity of I_Σ equals the projective dimension of $k[\Sigma^\vee]$, and vice versa.

2. Threshold graphs

We will use 3. in Theorem 1.7 to study threshold graphs.. We will denote a graph constructed in this way by $G = I_{1,1} \dots I_{1,i_1} S_1 I_{2,1} \dots I_{2,i_2} \dots S_2 I_{3,1} \dots I_{3,i_3} \dots S_{k-1} I_{k,1} \dots I_{k,i_k}$, where the I 's stand for adding an isolated vertex, and the S 's for adding a star vertex. We call this a presentation of G . We demand that $i_1 > 0$ to get uniqueness (one can start with a I or with an S), but all other i_j may be 0. To adjoin an I at the end corresponds to adding a new variable, which does not affect the Betti numbers, so when we are interested in Betti numbers, we can suppose that the presentation ends with an S . It is not hard to see that the maximal independent sets of G are $I_{1,1} \dots I_{k,i_k}$ (all I 's) and all $S_i I_{i+1,1} \dots I_{k,j_k}$ (each S followed by all I 's which come after S).

In a very interesting paper [4], Engström and Stamps consider Boij-Söderberg decomposition of rings with 2-linear resolution. Among other things they show that $k[G]$ has a 2-linear resolution if G is a threshold graph by using Theorem 1.6. They also show how the Betti numbers change when adding an I or an S , thereby giving a method to determine the Betti numbers for any threshold graph. For our purpose it is easier to use Theorem 1.7 to show 2-linearity and Theorem 1.8 to get the Hilbert series, and thus the Betti numbers. We give an example.

Example 2.1. Consider the threshold graph G with presentation $I_1 S_2 S_3 I_4 I_5 S_6$. The maximal independent sets are $\{I_1, I_4, I_5\}, \{S_2, I_4, I_5\}, \{S_3, I_4, I_5\}$, and $\{S_6\}$. We start with $\{I_1, I_4, I_5\}$ with series $1/(1-t)^3$ and attach $\{S_2, I_4, I_5\}$ with series $1/(1-t)^3$ in $\{I_4, I_5\}$ with series $1/(1-t)^2$. So far we get the series $2/(1-t)^3 - 1/(1-t)^2$. Then we attach $\{S_3, I_4, I_5\}$ with series $1/(1-t)^3$ in $\{I_4, I_5\}$ with series $1/(1-t)^2$ and get the series $3/(1-t)^3 - 2/(1-t)^2$. Finally we attach $\{S_6\}$ with series $1/(1-t)$ in the empty set with series 1. The final result is $k[G](t) = 3/(1-t)^3 - 2/(1-t)^2 + 1/(1-t) - 1 = (1 - 8t^2 + 15t^3 - 12t^4 + 5t^5 - t^6)/(1-t)^6$. Thus $\beta(k[G]) = (8, 15, 12, 5, 1)$.

The following picture illustrates the graph of $I_1 S_2 S_3 I_4 I_5 S_6$.

In the following theorem and corollaries we suppose that we have threshold graphs.

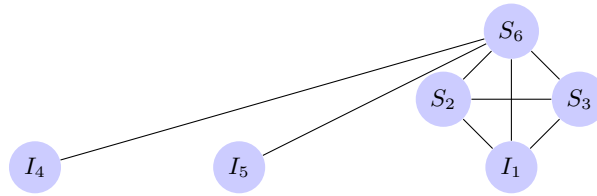


Figure 1. Graphs of $I_1S_2S_3I_4I_5S_6$.

Theorem 2.2. *The dimension of $k[G]$, equals the number of I 's in the presentation of G . Suppose that the presentation ends with an S . Then the depth of $k[G]$ is always 1. Both the dimension and the depth increases by one for each I we have after the last S in the presentation.*

Proof. We use Theorem 1.8. We have that $\max\{d_i + 1\}$ equals the number of I 's in the presentation of G , since the number of I 's in the presentation is at least one more than the number of I 's that follow an S , since the presentation starts with an I . Now $\min\{r_i + 1\} = 1$ if the presentation ends with an S . Adding an I at the end corresponds to adjoining a new variable which increases both dimension and depth by one. \square

Corollary 2.3. *We have that $k[G]$ is Cohen-Macaulay if and only if it has a presentation consisting of only I 's or is of the form I followed by a number of S 's, and possibly followed by a number of I 's. Thus $k[G]$ is Cohen-Macaulay if and only if G is either a graph of isolated vertices, or a complete graph and a (possibly empty) set of isolated vertices.*

Proof. If the presentations consists of only I 's, we have a polynomial ring. If not, we can suppose the presentation ends with an S by above. Then $\dim(k[G]) = \text{depth}(k[G]) = 1$ if and only if G has presentation $ISS \cdots S$. \square

Corollary 2.4. *Suppose $k[G]$ has presentation $I_1S_2S_3 \cdots S_n$. The Hilbert series of $k[G]$ is $(1 + (n - 1)t)/(1 - t)$. In this case $k[G]^\vee$, the Alexander dual of $k[G]$, has an $(n - 1)$ -linear resolution with nonzero Betti numbers $\beta_{0,0} = 1, \beta_{1,n-1} = n, \beta_{2,n} = n - 1$ and Hilbert series*

$$(1 - nt^{n-1} + (n - 1)t^n)/(1 - t)^n = (1 + 2t + 3t^2 + 4t^3 + \cdots + (n - 1)t^{n-2})/(1 - t)^{n-2}.$$

Proof. $I(G)$ has regularity 2 and $k[G]$ has projective dimension $n - 1$, so $I(G^\vee)$ has regularity $n - 1$ and $k[G]^\vee$ has projective dimension 2 according to Theorem 1.9. In case $k[G]$ is Cohen-Macaulay, $I(G^\vee)$ is generated by all n squarefree monomials of degree $n - 1$, so for $\beta_{i,j}(k[G]^\vee)$ we get $\beta_{1,n-1} = n$. Since $\beta_{0,0} = 1$, the resolution is linear, and the alternating sum $\beta_{0,0} - \beta_{1,n-1} + \beta_{2,n} = 0$, we get $\beta_{2,n} = n - 1$, and these are the only nonzero Betti numbers. The Hilbert series is thus $(1 - nt^{n-1} + (n - 1)t^n)/(1 - t)^n$. \square

We end this section with a list over threshold graphs with at most 5 vertices and presentation ending with an S .

Table 1. Threshold graphs with at most five vertices.

Type	β	dim	depth	f -vector
IS	(1)	1	1	(1,2)
IIS	(2,1)	2	1	(1,3,1)
ISS	(3,2)	1	1	(1,3)
IIIS	(3,3,1)	3	1	(1,4,3,1)
ISIS	(4,4,1)	2	1	(1,4,2)
IISS	(5,6,2)	2	1	(1,4,1)
ISSS	(6,8,3)	1	1	(1,4)
IIIS	(4,6,4,1)	4	1	(1,5,6,4,1)
SIIS	(5,7,4,1)	3	1	(1,5,6,2)
IISIS	(6,9,5,1)	3	1	(1,5,4,1)
IIISS	(7,12,8,2)	3	1	(1,5,3,1)
SISS	(7,11,6,1)	2	1	(1,5,3)
SISS	(8,14,9,2)	2	1	(1,5,2)
IISSS	(9,17,12,3)	2	1	(1,5,1)
SISSS	(10,20,15,4)	1	1	(1,5)

3. Neighborhood complexes for threshold graphs

Now we will determine which threshold graphs have Cohen-Macaulay neighborhood (or equivalently dominance) complexes.

Theorem 3.1. *The threshold graph G has a Cohen-Macaulay neighborhood complex, i.e., $k[N(G)]$ is Cohen-Macaulay, if and only if $k[G]$ is Cohen-Macaulay. In this case $k[N(G)]$ is a hyperplane section. Furthermore this occurs precisely when all facets in $N(G)$ have the same dimension.*

Proof. We can suppose that the presentation ends with an S , because I 's at the end have no neighbors, so they only contribute with new variables. If the presentation is $I_1 \cdots S_n$, then S_n has $n - 1$ neighbors, the maximal possible number. In the Stanley-Reisner ring $k[\Sigma]$ which is Cohen-Macaulay, all facets in Σ have the same dimension. Thus I_1 must have $n - 1$ neighbors if $N(G)$ is Cohen-Macaulay. This gives that there are only S 's after I_1 . In this case $N(G)$ is the $(n - 2)$ -dimensional skeleton of the $(n - 1)$ -dimensional simplex on $\{I_1, S_2, \dots, S_n\}$, which gives $k[N(G)] = k[x_1, \dots, x_n]/(x_1x_2 \cdots x_n)$. \square

We finish by giving a list of Hilbert series, dimension, depth of $k[N(G)]$ and f -vector of $N(G)$ for all threshold graphs with at most five vertices and presentation ending with an S .

Table 2. Hilbert series, dimension, depth of $k[N(G)]$ and f -vector of $N(G)$ for all threshold graphs.

Type	$k[N(G)](t)$	dim	depth	f -vector
IS	$(1+t)/(1-t)$	1	1	(1,2)
IIS	$(1+t-t^2)/(1-t)^2$	2	1	(1,3,1)
ISS	$(1+t+t^2)/(1-t)^2$	2	2	(1,3,3)
IIIS	$(1+t-2t^2+t^3)/(1-t)^3$	3	1	(1,4,3,1)
ISIS	$(1+t-t^3)/(1-t)^3$	3	2	(1,4,5,1)
IISS	$(1+t+t^2-t^3)/(1-t)^3$	3	2	(1,4,6,2)
ISSS	$(1+t+t^2+t^3)/(1-t)^3$	3	3	(1,4,6,4)
IIIS	$(1+t-3t^2+3t^3-t^4)/(1-t)^4$	4	1	(1,5,6,4,1)
ISIIS	$(1+t-t^2-t^3+t^4)/(1-t)^4$	4	2	(1,5,8,4,1)
IISIS	$(1+t-2t^3+t^4)/(1-t)^3$	4	2	(1,5,9,5,1)
IIISS	$(1+t+t^2-2t^3+t^4)/(1-t)^4$	4	2	(1,5,10,7,2)
ISSIS	$(1+t-t^4)/(1-t)^4$	4	3	(1,5,9,7,1)
ISISS	$(1+t+t^2-t^4)/(1-t)^4$	4	3	(1,5,10,9,2)
IISSS	$(1+t+t^2+t^3-t^4)/(1-t)^4$	4	3	(1,5,10,9,3)
ISSSS	$(1+t+t^2+t^3)/(1-t)^4$	4	4	(1,5,10,10,5)

4. Conclusion

Edge rings of threshold graphs were shown to have 2-linear resolution by Engström and Stamps. They used that threshold graphs are co-cordal. We determined when they are Cohen-Macaulay by instead showing that they have Stanley-Reisner rings which are fat forests. We also determined when the neighborhood complex of a threshold graph is Cohen-Macaulay.

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